New perspectives on protein flexibility
High dimensional volumes and DoS
Move sets for polypeptide chains
Comparing energy landscapes

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3IA Côte d'Azur Axis 3

## Overall perspective

$\triangleright$ When is a well-posed computer science/modeling problem solved?

- Intrinsic difficulty understood
- (Almost) Optimal algorithms available
$\triangleright$ Strategy:
- Identify computationally tractable problems
- Approximability is the issue, not NP-hardness
- Develop efficient algorithms
- Bias on the geometric/combinatorial side
- Develop the corresponding software
- Software: large research instrument
$\triangleright$ Structure, thermodynamics, kinetics: will these problems get solved ?

New perspectives on protein flexibility

Volumes of polytopes

```
Tripeptide Loop Closure (TLC)
    TLC: background
    TLC steric constraints
    Loop sampling
```


## Comparing energy lanscapes

## Landscapes and thermodynamics

Density of states and partition functions
Dialanine


- Potential energy:

$$
V_{\text {total }}=V_{\text {bonded }}+\left(V_{\text {vdw }}+V_{\text {electro }}\right)
$$

$\triangleright$ Potential energy landscape:

$\triangleright$ Density of states (DoS) for $A \subset X$ :

- For any $v_{0}<v$ :

$$
G\left(\left[v_{0}, v\right]\right)=\int_{A} 1_{\left[v_{0}, v\right]}(V(x)) d x
$$

$\triangleright$ Partition function for $A \subset X$ from DoS:

$$
Z_{A}(T)=\int_{A} e^{-\beta v} d G(v)
$$

$\triangleright \mathrm{Nb}$ : DoS calculation: volume calculation in phase space

## Polytope volume calculations

$\triangleright$ Problem statement: design effective algorithms to estimate the volume of high dimensional polytopes (dim. $\in[100 \ldots$. 1000])

$\triangleright$ Unless $\mathrm{P}=\mathrm{NP}$ : no polynomial time algorithm with approx factor $(c d / \log d)^{d}$
$\triangleright$ State-of-the-art: multi-phase Monte Carlo methods embarking

- Rounding procedures to put the polytope in isotropic position
- Random walks: ball-walk, hit-and-run, billiard walk
- Mixing times analysis - and heuristics for early stops

```
\trianglerightRef: Cousins and Vempala, Math. Prog. Comp., }201
Ref: Chalkis, Emiris, Fisikopoulos, arXiv:1905.05494, 2019
\trianglerightRef: Chevallier et al, J. Computational Geometry, 2022
Ref: Chevallier et al, AISTATS, 2022
```


## Volume of polytopes: hardness, randomized algorithms


$\triangleright$ Hardness: no polynomial time algorithm with approx factor $(c d / \log d)^{d}-$ unless $P=N P$
$\triangleright \varepsilon$-approximation of the volume: for any parameter $\varepsilon>0$, a number $V$

$$
(1-\varepsilon) \operatorname{Vol}(K) \leq V \leq(1+\varepsilon) \operatorname{Vol}(K)
$$

$\triangleright(\varepsilon, \delta)$-approximation algorithm: algorithm returning an $\varepsilon$-approximation with a probability at least $1-\delta$.
$\triangleright$ Complexity, the $O^{\star}(n)$ otation:

- $O\left(d^{4}\right)$ : upper bound as a function of the dimension $d$
$-O^{\star}\left(d^{4}\right)$ : term in $\log d, \varepsilon, \delta$ removed; focus on the dimension solely
$\triangleright$ Ref: Cousins, Vempala, SIAM J. Comp., 2018


## Random walk: hit-and-run

$\triangleright$ Goal: sample point in $K$ according to a prescribed density $f$
$\triangleright$ (Random-direction) hit-and-run: random point $x_{W}$ after $W$ steps

$\triangleright$ Iteratively:

- pick a random vector
- move to random point on the chord $I \cap K$, chosen from the distribution induced by $f$ on $I$
$\triangleright$ Comments:
- risk of being trapped near a vertex
- large $W$ helps forgetting the origin $x_{0}$
$\triangleright$ Thm (Berbee et al) The limit distribution induced by HR is uniform in $K$.
$\triangleright$ Thm (Vempala et al) HR can be modified to sample an isotropic Gaussian (restricted to K).
$\triangleright$ Thm (Lovász) Let $r$ and $R$ denote the radii of the largest inscribed and circumscribed balls for $K$. One sample generation: $O^{\star}\left(d^{3}\right)$.
$\triangleright$ NB: precise statement in terms of total variation distance omitted
$\triangleright$ Ref: Berbee et al, Math. Prog., 1987
$\triangleright$ Ref: Lovász, Math. Prog. Ser. A, 1999


## Randomized algorithms: complexity

$\triangleright$ Volume estimated using a sequence of isotropic Gaussians:

$$
\begin{equation*}
\operatorname{Vol}(K)=\int_{K} f_{0}(x) d x \frac{\int_{K} f_{1}(x) d x}{\int_{K} f_{0}(x) d x} \cdots \frac{\int_{K} d x}{\int_{K} f_{m-1}(x) d x} \equiv \int_{K} f_{0}(x) d x \prod_{i=1, \ldots, m} R_{i} \tag{1}
\end{equation*}
$$

$\triangleright$ Cooling schedule i.e. sequence of Gaussians $f_{0}, \ldots, f_{m}$ :

- $f_{0}$ : sharply peaked in $K$
- $f_{m}$ : uniform distribution i.e. $a_{m}=0$
$\triangleright$ Thm. For a convex body $K$ given by a membership oracle, and such that
$B \subset K \subset R B$, an $(\varepsilon, \delta)$ - approximation can be obtained in time

$$
\begin{equation*}
O\left(\frac{d^{4}}{\varepsilon^{2}} \log ^{9} \frac{n}{\varepsilon \delta}+d^{4} \log ^{8} \frac{n}{\delta} \log R\right)=O^{\star}\left(d^{4}\right) \tag{2}
\end{equation*}
$$

$\triangleright$ Ref: Lovász, Vempala, J Comp. Syst. Sciences, 2006
$\triangleright$ Ref: Cousins, Vempala, SIAM J. Comp., 2018

## A practical algorithm: outline

$\triangleright$ Method:

- multi-phase Monte-Carlo using $m=O(\sqrt{d})$ logconcave functions $\left\{f_{0}, \ldots, f_{m-1}\right\}$,
- $f_{i}(x) \propto e^{-a_{i}^{T} x}$ or $f_{i}(x) \propto \exp \left(-a_{i}\|x\|^{2}\right)$
- At each step: estimate $r_{k} \approx \int_{K} f_{k}(x) d x / \int_{K} f_{k-1}(x) d x$

```
\(\operatorname{Volume}(K, \varepsilon)\) : Convex body \(K\), error parameter \(\varepsilon\).
\(-T=\operatorname{Round}\left(\right.\) body: \(K\), steps: \(8 n^{3}\) ), set \(K^{\prime}=T \cdot K\).
- \(\left\{a_{0}, \ldots, a_{m}\right\}=\) GetAnnealingSchedule(body: \(K^{\prime}\) ).
- Set \(x\) to be random point from \(f_{0} \cap K^{\prime}, \varepsilon^{\prime}=\varepsilon / \sqrt{m}\).
- For \(i=1, \ldots, m\),
    - Set \(k=0, x_{0}=x\), converged \(=\) false, \(W=4 n^{2}+500\).
    - While converged \(=\) false,
            - \(k=k+1\).
            - \(x_{k}=\) HitAndRun(body: \(K\), target distribution: \(f_{i-1}\), current point: \(x_{k-1}\) ).
            - Set
                \(r_{k}=\frac{1}{k} \sum_{j=1}^{k} \frac{f_{i}\left(x_{j}\right)}{f_{i-1}\left(x_{j}\right)}\).
            - Set \(W_{\max }=\max \left\{r_{k-W+1}, \ldots, r_{k}\right\}\) and \(W_{\min }=\min \left\{r_{k-W+1}, \ldots, r_{k}\right\}\).
            - If \(W_{\text {max }}-W_{\text {min }} \leq \varepsilon^{\prime} / 2 \cdot W_{\text {max }} \rightarrow\) converged \(=\) true.
    - Set \(R_{i}=r_{k}, x=x_{k}\).
- Return volume \(=|T| \cdot\left(\pi / a_{0}\right)^{n / 2} \cdot R_{1} \ldots R_{m}\).
```

$\triangleright$ Ref: Cousins and Vempala, Math. Prog. Comp., 2016

## Piecewise deterministic Markov processes (PDMP)

## the non-reversible Bouncy Particle Sampler (BPS)

$\triangleright$ Notations: state space (position, velocity): $z=(x, v) \in E=\mathbb{R}^{d} \times \mathbb{R}^{d}$.
$\triangleright$ PDMP $z_{t}$ : a continuous time Markov process defined by:

1. a deterministic flow $\phi_{t}(z)$,
2. function determining the length of steps: jump kernel $\lambda(z)$
3. a jump kernel in phase $(x, v)$
space: $q(\cdot \mid z)$
$\triangleright$ BPS: PDMP to sample a distribution $\pi(x)$ in $\mathbb{R}^{d}$ using piecewise linear trajectories bouncing on high energy level set surfaces
4. Linear trajectories: $\phi_{t}(x, v)=(x+t v, v)$,
5. Arrival time of 1 D inhomogeneous Poisson process of intensity $\lambda(x, v)=\max \left(0,-\left\langle\nabla_{x}(\log \pi)(x), v\right\rangle\right)$,
6. $q(\cdot \mid z)$ : reflection w.r.t. the gradient of the potential:

$$
\begin{equation*}
\left(x, v^{\prime}\right)=\left(x, v-2 \frac{\left\langle v, \nabla_{x}(\log \pi)(x)\right\rangle}{\left\|\nabla_{x}(\log \pi)(x)\right\|^{2}} \nabla_{x}(\log \pi)(x)\right) \tag{3}
\end{equation*}
$$

4. +Refresh of velocity to ensure ergodicity
$\triangleright$ Ref: Doucet et al, Stats. and probability letters, 136, 2018

## Extension: BPS on a bounded domain - a polytope

$\triangleright$ Example BPS trajectory in the 2d cube $[-1,1]^{2}$ :
$\triangleright$ Three types of events:

- PDMP events: as usual
- Reflexions on the boundary

$$
\begin{equation*}
v^{\prime}=v-2 \frac{\langle n, v\rangle}{\|n\|^{2}} n \tag{4}
\end{equation*}
$$

- Refresh events: velocity resampled from isotropic normal distribution
$\triangleright$ Numerics: lazy update of linear algebra operations


Blue: PDMP jump events, Red: reflections on the boundary, Green:
refresh events
$\mathrm{Nb}: \pi(x)$ : Gaussian of variance
$\sigma=1$.

## PDMP to compute volumes of polytopes: experiments

$\triangleright$ Complexity: $\quad C=O\left(d^{c}\right)$, dimension up to $d=250$
$\triangleright$ Protocol: find the smallest number of samples so that the estimated volume is within err\% from the exact value


$\triangleright$ Linear regression in $\log \log$ scale for the three polytopes:

|  | Time |  | Num. samples |  |
| :---: | :---: | :---: | :---: | :---: |
| model | slope | $R^{2}$ | slope | $R^{2}$ |
| cube | 3.77 | 0.96 | 1.94 | 0.88 |
| $\Delta_{\text {iso }}$ | 3.52 | 1.00 | 1.72 | 0.99 |
| $\Delta_{\text {std }}$ | 3.18 | 0.99 | 1.37 | 0.96 |

## Computing volumes and DoS: outlook

- Polytopes: very efficient algorithms, provably correct
- Beyond polytopes: three classes of questions
- Designing cooling schedules
- Mixing times of RW - related to the conductance of the Markov chains i.e. narrow passages
- Sample generation - beyond line-segments


## Bibliography: volumes

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J. of Computational Physics, 410(1):1-19, 2020.

## New perspectives on protein flexibility

## Volumes of polytopes

Tripeptide Loop Closure (TLC)
TLC: background
TLC steric constraints
Loop sampling

## Comparing energy lanscapes

## Geometric models: Cartesian and internal coordinates

$\triangleright$ Cartesian versus internal coordinates: $\left\{x_{i} y_{i} z_{i}\right\}_{i}$ versus $\left\{d_{i j}, \theta_{i j k}, \sigma_{i j k l}\right\}$
$\triangleright$ Bond length and valence angle
(A)

(B)

$\triangleright$ Dihedral angles

(B)

$\triangle$ Protein backbone


Ramachandran diagram, per a.a. type:
$\triangleright$ Side chain: 20 natural amino acids Exple: Lysine, 4 dihedral angles

- bivariate distribution for $(\phi, \psi)$


LYS

## Softness of Internal coordinates -force constants from снавмм 36



Bonds: $\delta d_{i j} \sim .2 \AA: \Delta V \sim 20 \mathrm{kcal} / \mathrm{mol}$


Torsion angles: $\Delta V \sim 3-4 \mathrm{kcal} / \mathrm{mol}$


Valence angles: $\delta \theta_{i j} \sim 10^{\circ}: \Delta V \sim 20 \mathrm{kcal} / \mathrm{mol}$


Quadrature vs importance sampling (Frenkel and Smit, 2002)
$\Rightarrow$ Dihedral angles are indeed soft coordinates

## The Ramachandran diagrams

$\triangleright$ Ramachandran diagrams and populated regions


- Main regions: $\alpha L, \alpha R, \beta S, \beta P$
- Three prototypical diagrams
- Glycine
- Proline
- Others - e.g. Aspartic acid
$\triangleright$ Distance constraints and the Ramachandran tetrahedron
$C 1: C_{\beta}-O_{i-1} \quad C 2: C_{\beta}-O+C_{\beta} N_{i+1}$ C3: $O_{i-1}-O+O_{i-1} N_{i+1}$

$\triangleright$ Ref: Stereochemistry of polypeptide chain configurations, JMB, 1963;
Ramachandran et al
$\triangleright$ Ref: Revisiting the Ramachandran plot, Protein Science, 2003; Ho et al


## The Tripeptide loop closure - TLC

$\triangleright$ TLC: for 3 amino acids, fix all internal coordinates BUT the $\left(\phi_{i}, \psi_{i}\right)_{i=1,2,3}$ angles

$\triangleright$ Theorem: at most 16 solutions


3 consecutive a.a.
$\Rightarrow$ Find all possible values
$\left(\phi_{i}, \psi_{i}\right)_{i=1,2,3}$ compatible with the fixed internal coordinates


3 a.a. sandwiching SSE-CDRs
$\triangleright$ Ref: Gō and Scheraga, Macromolecules, 1970
$\triangleright$ Ref: Coutsias et al, J. Comp. Chem., 2004

## TLC model: from six to three angles

$\triangleright$ Motions of the 3 rigid bodies: 6 angles



Nb : indices $\bmod (3)$, e.g., $\sigma_{0}=\sigma_{3}$
$\triangleright \ldots$ which are actually three

$$
\begin{equation*}
\sigma_{i}=\tau_{i}+\delta_{i} \tag{5}
\end{equation*}
$$


$\delta_{i}=\angle \operatorname{Plane}\left(C_{\alpha ; i} C_{\alpha ; i+1} C_{i}\right)$, Plane $\left(C_{\alpha ; i} C_{\alpha ; i+1} N_{i+1}\right)$
$\triangleright$ Key ingredients of TLC:

- Initially: six dihedral angles $\{(\phi, \psi)\}_{\{i=1,2,3\}}$
- Then: three pairs $\left\{\delta_{i}, \tau_{i}\right\}$
- Finally: three angles $\tau_{i}$
$\triangleright$ The valence angle constraints: the $\theta_{i}$ angles at the $C_{\alpha ; i}$ s must remain constant.
$\Rightarrow$ It is the coupling introduced by the $\theta_{i}$ angles onto the rotation angles $\tau_{i}$ yields a degree 16 polynomial.
$\triangleright$ Ref: Coutsias et al, 2004


## TLC with moving legs and embeddable tripeptides

$\triangleright$ Geometric model:

- Tripeptide such that: left leg $N_{i} C_{\alpha ; i}$ fixed, right leg $C_{\alpha ; i+2} C_{i+2}$ free to move
- Six dihdedral angles $\left\{\phi_{i}, \psi_{i}\right\}$ free
$\triangleright$ Question: provide necessary conditions on the position of the first and last segment-the legs, for the Tripeptide Loop Closure (TLC) algorithm to hold solutions. $\triangleright \mathrm{Nb}$ : the relative position of legs suffices; in that case, position + orientation of $C_{\alpha ; i+2} C_{i+2}$ yields a 5 -dim search space.



## TLC: necessary conditions on the existence of solutions

$\triangleright$ TLC problem for a tripeptide - say $T_{k}$ : degree 16 polynomial parameterized by 12 angles defining the space $\mathcal{A}_{k}=\left\{\alpha_{k, i}, \eta_{k, i}, \xi_{k, i-1}, \delta_{k, i-1}\right\}, i \in\{1,2,3\}$.

$\triangleright$ Contribution: necessary conditions for TLC to admit solutions

- Based on the 12 angles in $\mathcal{A}_{k}$
- Defined by 24 hyper-surfaces in $\mathcal{A}_{k}$
- These hyper-surfaces: curved walls for Hit-and-Run
$\triangleright$ Ref: O'Donnell, Cazals; J. Comp. Chem., 2023


## Polytope volume calculations

$\triangleright$ Problem statement: design effective algorithms to estimate the volume of high dimensional polytopes (dim. $\in[100 \ldots$. 1000])

$\triangleright$ Unless $\mathrm{P}=\mathrm{NP}$ : no polynomial time algorithm with approx factor $(c d / \log d)^{d}$
$\triangleright$ State-of-the-art: multi-phase Monte Carlo methods embarking

- Rounding procedures to put the polytope in isotropic position
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Ref: Chevallier et al, AISTATS, 2022
```


## Global geometric model

$\triangleright$ Loop studied $L$ : $M=3 \times m$ amino, $m$ tripeptides: $L=T_{1}, \ldots, T_{m}$
$\triangleright$ Loop decomposition: rigid peptide bodies and their complements

$$
\begin{equation*}
L=P_{0} T_{1}^{\prime} P_{1} \ldots P_{k-1} T_{k}^{\prime} P_{k} \ldots P_{m-1} T_{m}^{\prime} P_{m} \tag{6}
\end{equation*}
$$


$\triangleright$ Parametric space:

- For one peptide body: $S E(3)=S O(3) \times \mathbb{R}^{3}$
- For one tripeptide: solution space of TLC. .. except that
- The angular parameterization of $\operatorname{TLC}\{\alpha, \xi, \eta, \delta\}$ : depends on $S E(3) \times S E(3)$ since the left and right legs come from $P_{i-1}$ and $P_{i-1}$


## Loop sampling: spaces involved and solution sketch

$\triangleright$ Loop decomposition into: rigid peptide bodies and tripeptides cores


$$
\begin{aligned}
L= & P_{0} T_{1}^{\prime} P_{1} \ldots \\
& P_{k} T_{k+1}^{\prime} P_{k+1} \ldots \\
& P_{m-1} T_{m}^{\prime} P_{m} .
\end{aligned}
$$

$\triangleright$ Random sampling of loop conformations using Hit-and-Run:


- Aim: perform rejection sampling in a region $\mathcal{V}$ containing all valid loop geometries.
- How: with Hit-and-Run in a domain characterizing necessary conditions - cf validity intervals


## Loop sampling: spaces involved and solution sketch

$\triangleright$ Global parameterization of the conformational space of the loop: based on rigid bodies associated with peptide bonds

- $\mathcal{M}$ : motion space for the $m-1$ peptide bodies, essentially $(S E(3))^{m-1}$
- A: 12m-dimensional angular space coding the geometry of tripeptides
- $\mathcal{V}$ : domain bounded by hyper-surfaces corresponding to Validity Constraints Necessary Constraints for TLC to admit solutions
- $\mathcal{S}$ : the fertile space, where TLC admits one solution for each tripeptide
- $\mathcal{F}$ : clash free solutions in $\mathcal{S}$ for $\left\{N, C_{\alpha}, C, O, C_{\beta}\right\}$ pairs
$\triangleright$ Number of solutions: $\prod_{i}($ num solutions tripeptide $i)$



## Validity domain for the whole chain $L$ with $m$ tripeptides

$\triangleright$ Angles $\tau$ : $3 m$ angles $\tau$ (3 for each tripeptide)
$\triangleright$ Recap per angle $\tau$ :

- For one angle: at most 4 Depth One Validity Intervals (DOVI)
- For each DOVI: 2 sub-manifolds of $\mathcal{A}_{k}$ defined by the previous equations; yields (at most) 8 sub-manifolds in $\mathcal{A}_{k}$.
$\triangleright$ For one tripeptide: $3 \tau$ angles $\Rightarrow 24$ constraint surfaces in the 12 dimensional angular space $\mathcal{A}_{k}$.
$\triangleright$ For the whole loop: total of $24 m$ constraint surfaces.



## Algorithms and parameters

$\triangleright$ Unmixed loop sampler ULS ${ }_{\text {One|All; } N_{E S}}^{N_{V} ; N_{O R}}\left[p_{0}\right]$ :

- One $\mid$ All a flag indicating how many solutions are retained at each embedding step,
- $N_{E S}$ the number of embedding steps,
- $N_{V}$ the number of random trajectories followed in motion space,
- $N_{O R}$ the output rate (the number of steps in-between the ones where conformations get harvested),
- $p_{0}$ : the starting configuration.
$\triangleright$ Mixed loop sampler $\mathbb{M L L} \mathbb{S}_{\text {One } \mid \text { All } ; N_{E S}}^{N_{V} ; N_{O R}}\left[p_{0}\right]$ : every other step, the loop is shifted by 1 or 2 units to also sample the peptide bodies.


## Loops sampling: $\phi, \psi$ and $\omega$

$\triangleright$ Typical values of the torsion angle $\omega$ :

- SSE?
- loops?


## Loops sampling: $\phi, \psi$ and $\omega$

$\triangleright$ Typical values of the torsion angle $\omega$ :

- SSE? $\pi \pm 2-3^{\circ}$
- loops? $\pi \pm 15^{\circ}$



## Illustration: CDR-H3-HIV, 30 amino acids

$\triangleright$ System:

- The loop is a complementarity-determining region (CDR-H3) from PG16, an antibody with neutralization effect on HIV-1.
- pdbid: 3mme, chain A; residues: 93-100, 100A-100T, 101, 102.


Conformations generated by algorithm $\mathbb{M L} \mathbb{S}^{1 ; 1}$ the 30 a.a. long CDR3. (B,C) Side/top view of 250 conformations.
$\triangleright$ Generation speed: $\sim 10$ conformations per second

Results: sampling and study of fluctuations


Backbone RMSF (36 atoms) for the 12 amino acid long loop PTPN9-MEG2.

## Bibliography: backbone move sets

國 T. O'Donnell, C.H. Robert, and F. Cazals.
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J. Comp. Chem., 2023.

## Outlook

$\triangleright$ Key features:

- First global parametric model of protein loops amenable to effective sampling strategies a-la Hit-and-Run
- Results: on par or better with state-of-the-art methods
- Atomic fluctuations along the loop
- Mutual reachability for existing conformations
- Insights on the intrinsic difficulty of the problem-via random walks and curved polytopes
$\triangleright$ Open problems:
- Uniformity of sampling (Theorem)
- Connexion to micro-canonical ensembles and densities of states
- Sampling with side chains

New perspectives on protein flexibility

## Volumes of polytopes



Comparing energy lanscapes

## Comparing (Sampled) Energy Landscapes: Motivation

$\triangleright$ Comparing (sampled) landscapes:

- Assessing the coherence of two force2 fields for a given system (atomic, CG)
- Comparing two related systems: e.g. wild type/mutated proteins
- Comparing two simulations: different initial conditions and/or algorithms

$\triangleright$ Idea: find a mapping between basins considering
- the similarity between the native states (one per basin)
- the coherence between the volumes of the basins (their probabilities)
- the connectivity between basins
$\triangleright$ Terminology: sampled (potential) energy landscape:
- portion revealed by a simulation
- given: minima, transitions between them, volumes of basins


## Comparing Sets of Local Minima using a

Minimum Oriented Spanning Forest (MSF): method
$\triangleright$ Given two sets of local minima and a distance metric to compare them: each local minimum chooses its nearest neighbor cf One-sided Hausdorff distance


NB: local minima

- all those discovered during exploration
- persistent ones only (remove ruggedness)
$\triangleright$ Statistics:
- ave. weight of edges from the first landscape to the second one: $\bar{w}_{1 \rightarrow 2}^{M S F}$
- ave. weight of edges from the second landscape to the first one: $\bar{w}_{2 \rightarrow 1}^{\text {MSF }}$
$\triangleright$ Remarks:
- can be combined with topological persistence
- algorithm, of MST: Borůvka/ distributed Kruskal


## Comparisons without Connectivity Constraints:

## the Earth Mover Distance yields a Linear Program

$\triangleright$ Consider two landscapes: $\mathrm{PEL}_{s}$ with $n_{s}$ basins, $\mathrm{PEL}_{d}$ with $n_{d}$ basins
(a)


(b)

$\triangleright$ Problem Earth-Mover-Distance (EMD):
find the transport plan of minimum cost, i.e. solution of the following linear program

$$
L P \begin{cases}\text { Cost: } \operatorname{Min} \sum_{i=1, \ldots, n_{s}, j=1, \ldots, n_{d}} f_{i j} \times d_{\mathcal{C}}\left(s_{i}, d_{j}\right) & \\ \sum_{i=1, \ldots, n_{s}} f_{i j}=w_{j}^{(d)} & \forall j \in 1, \ldots, n_{d} \\ \sum_{j=1, \ldots, n_{d}} f_{i j} \leq w_{i}^{(s)} & \forall i \in 1, \ldots, n_{s} \\ f_{i j} \geq 0 & \forall i \in 1, \ldots, n_{s}, \forall j \in 1, \ldots, n_{d}\end{cases}
$$

$\triangleright$ Property: in OPT, the number of edges carrying flow is $O\left(n_{s}+n_{d}-1\right)$
$\triangleright$ Pros and cons:

- Information used: location of minima, weight of basins
- Linear program: solved in polynomial time
- Connectivity information not used
$\triangleright$ Ref: Chvátal, Linear programming, 1983; Rubner, Tomasi, Guibas, IJCV, 2000


## Comparisons with Connectivity Constraints

$\triangleright$ Earth Mover Distance: may violate the connectivity constraints
Landscape: source

$\triangleright$ Def: Transport plan with connectivity constraints: every connected subgraph of $\mathrm{PEL}_{s}$ exports towards a connected subgraph of $\mathrm{P} E L_{d}$

* There may exist an exponential number of connected subgraphs
$\triangleright$ Problem EMD-CCC: maximum flow under constraints of
\{maximum cost, connectivity constraints (and transport plan size $M$ ) \}
$\triangleright$ Complexity results
- Decision versions of EMD-CC and EMD-CCC: NP-complete
- Optimization version of EMD-CC is not in APX

If $P \neq N P$ : no polynomial algorithm with constant approx factor
$\triangleright$ Algorithm Alg-EMD-CCC-G

- Greedy polynomial algorithm producing solutions i.e.
respecting the connectivity constraints and the max cost.
Complexity: $O\left(n^{3} m^{2}\right)$, with $n$ and $m$ the num. vertices of the graphs


## Bibliography: comparing landscapes

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## The Structural Bioinformatics Library



Why adopt the SBL ?

## For Biologists:




For Developers:

- IrnatConprithn tootbra



Reference:

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$\triangleright$ Pointers:
- Frontpage
- Applications
- Online doc
$\triangleright$ Upates
- Conda channels for linux and macos
- Online demos for applications
- Next: plugins for VMD and pymol


## Acknowledgments

$\triangleright$ Collaborators

- Dorian Mazauric, Timothee Of Onnell (Inria)
- Augustin Chevallier (Univ. Lancaster)

Edoardo Sarti and Côme Le Breton (Inria)
7) Charles Robert (CNRS IPBC)

- Juan Cortés (CNRS LAAS)


An David Wales, (Cambridge University)
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